#### 1. Introduction and Overview

- Switching and Transmission
- Multiplexing and Concentration
- Timescales of Information Transfer

#### 2. Circuit Switch Design Principles

- Space-domain Circuit Switching
- i. Nonblocking Properties, Complexity
- ii. Clos, Benes and Cantor Switches
- Time-domain Switching

#### 3. Fundamental Principles of Packet Switch Design

- Packet Contention
- Interconnection Networks (Banyan Networks)
- Sorting Networks
- Nonblocking and Self-Routing Properties of Clos Networks

#### **4. Advanced Packet Switching Principles**

- Performances of Simple Switches
- Look-Ahead Contention Resolution
- Speedup Principle
- Channel-Grouping Principle
- Knockout Principle
- Replication Principle
- o Dilation Principle

#### **5. Advanced Switch Design Principles**

- Switch Design Principles Based on Deflection Routing
- i. Tandem-Banyan Network
- ii. Shuffle-Exchange Network
- iii. Feedback Shuffle-Exchange Network
- iv. Feedback Bidirectional Shuffle-Exchange Network
- v. Dual Shuffle-Exchange Network
- Switching by Memory I/O

#### 6. Switching Principles for Multicast, Multirate, and Multimedia Services

- Multicast Switching
- i. Multicasting Based on Nonblocking Copy Networks
- ii. Performance Improvement of Copy Networks
- iii. Multicasting Algorithm for Arbitrary Network Topologies
- iv. Nonblocking Copy Networks Based on Broadcast Clos Networks
- Path Switching
- i. Basic Concept of Path Switching
- ii. Capacity and Route Assignments for Multirate Traffic
- iii. Trade-Off Between Performance and Complexity
- iv. Multicasting in Path Switching

#### 7. Packet Switching and Information Transmission

- Duality of Switching and Transmission
- Parallel Characteristics of Contention and Noise
- i. Pseudo Signal-to-Noise Ratio of Packet Switch
- ii. Clos Network with Random Routing as a Noisy Channel
- Clos Network with Deflection Routing
- Cascaded Clos Network
- ii. Analysis of Deflection Clos Network
- Route Assignments and Error-Correcting Codes
- i. Complete Matching in Bipartite Graphs
- ii. Graphical Codes
- iii. Route Assignments of Benes Network
- Clos Network as Noiseless Channel-Path Switch
- i. Capacity Allocation
- ii. Capacity Matrix Decomposition
- Scheduling and Source Coding
- i. Smoothness of Scheduling
- ii. Comparison of Scheduling Algorithms
- iii. Two-Dimensional Scheduling

# Merging of computer and communications technologies

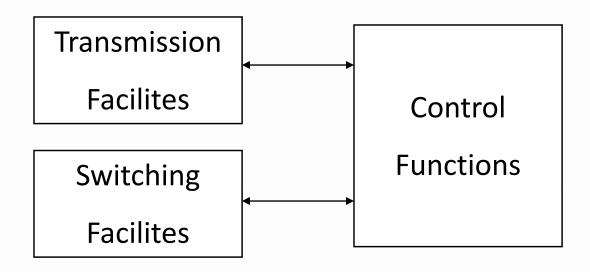
#### 2. Traditional Computer Networks

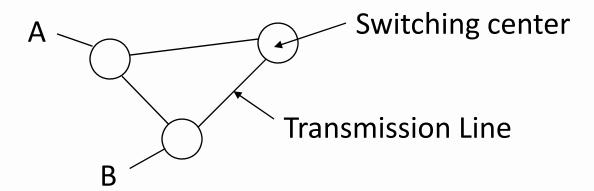
- Services non real time
- data rate not too high
- service quality not guaranteed

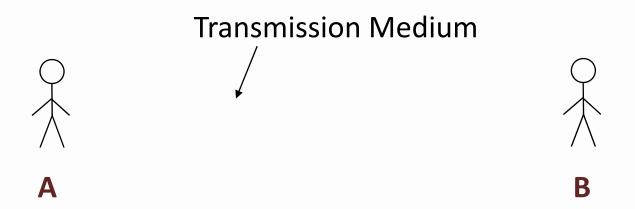
#### 3. Telephone Networks

- o Real time
- low data rate (64 kbps)
- Service quality guaranteed

- 4. Computer Networks on Telephone Networks
- o Modems
- leased line
- 5. Multimedia services
- 6. Integrated Communications Networks
- 7. Broadband
- 8. Integration
- 9. Support of unforeseen services



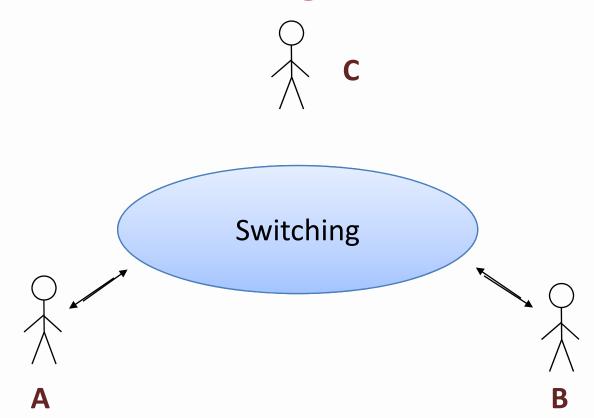




When there are only two users, information from A is by default destined for B, and vice versa

For a two-user network, switching is not required

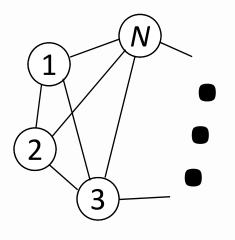
Fig. 1.1. A two-user network; switching is not required



Information from A may be destined for B or C

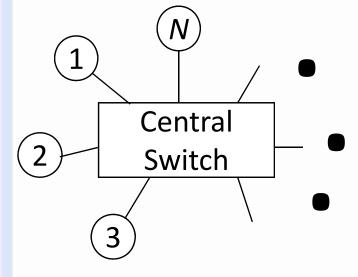
For a three-user network, switching is required

Fig. 1.2. A three-user network; switching is required



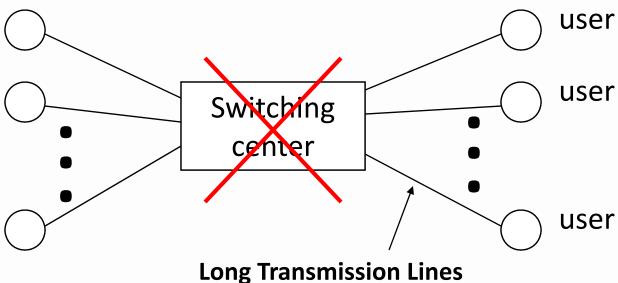
Switching is performed at user's location by selecting one of the links for reception

# of bi-directional links = N(N-1)/2



Switching is performed by a central switch

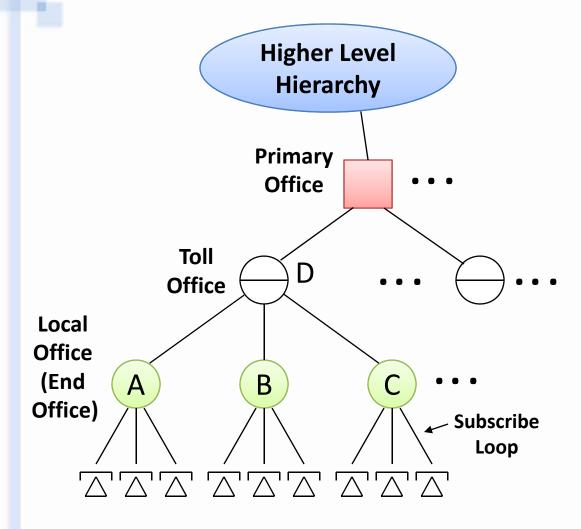
Fig. 1.3. N-user networks with switching performed (a) at users' location (b) by a central switch



In case the totally centralized switch is not attractive, no communication can be processed between any users

An example in which totally centralized switching is not attractive

#### **Transmission Hierarchy**



- i.) Decrease of Traffic volume as going up in hierarchy
- ii.) Resources at upper hierarchy shared by larger population
- iii.) Hierarchical routing

Fig. 1.4. Telephone network hierarchy

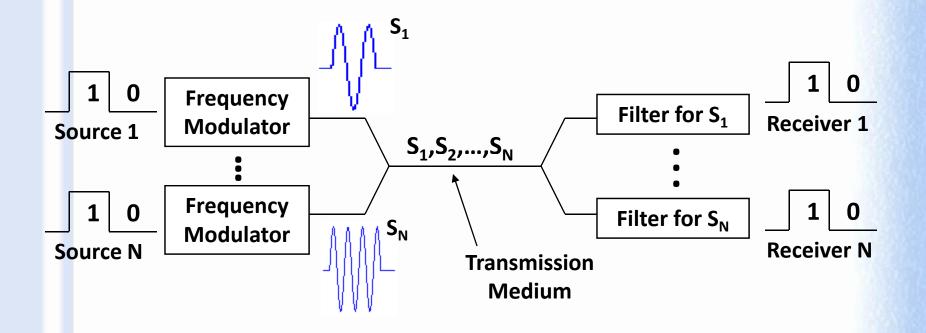


Fig. 1.5. Frequency-division multiplexing

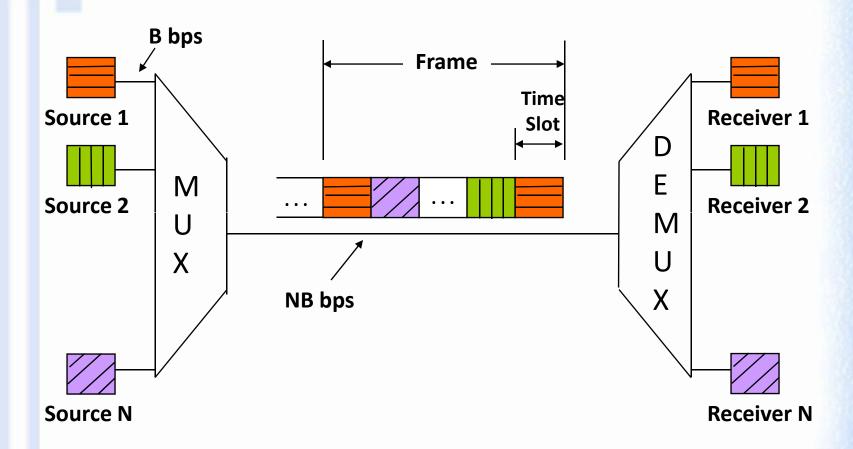
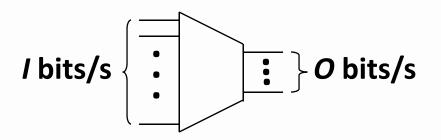


Fig. 1.6. Time-division multiplexing

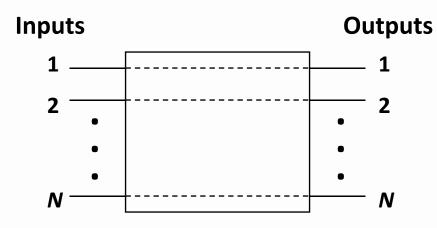
#### Difference between Multiplexing and concentration:



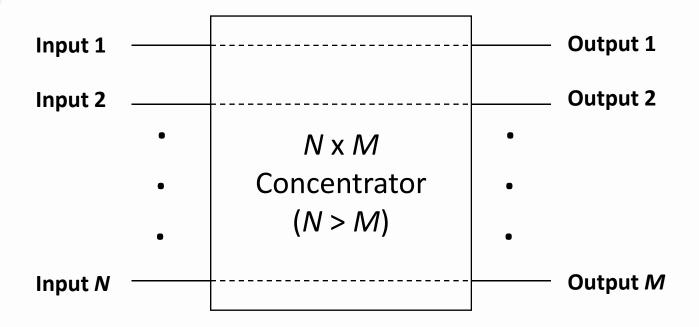
I = O in mux. system

*I > O* in concentrator

#### Difference between Multiplexing and Switching:

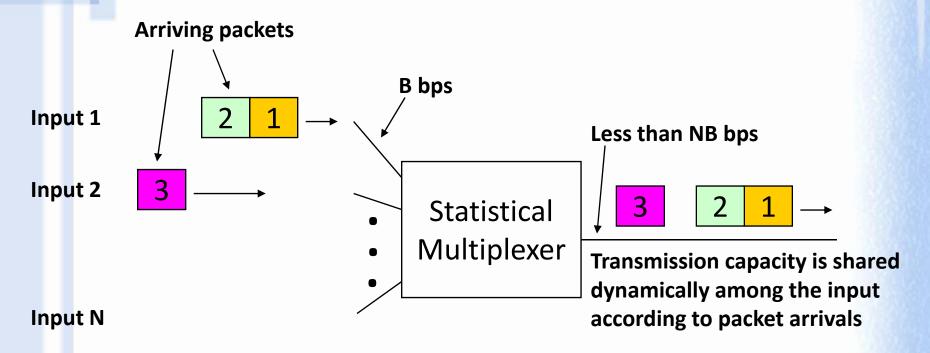


Connectivity is static in mux, system but dynamic in switching system



An active input is assigned to one of the outputs. It does not matter which output is assigned.

Fig. 1.7. An N x M concentrator



- o Combines multiplexing and concentration
- Packet header indicating destination
- Switching involved
- Connectivity changes with time

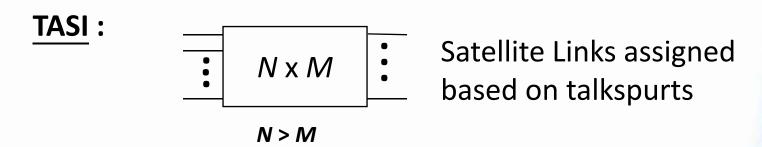
Fig. 1.8. A Statistical multiplexer

#### Sessions, circuits (concentration):

- Session needs to be set up before communication
- Exclusive and non-exclusive dedication of resources during connection
- Virtual circuits (not connection oriented)

#### Messages (Statistical multiplexing):

- variable packets
- o talk spurts



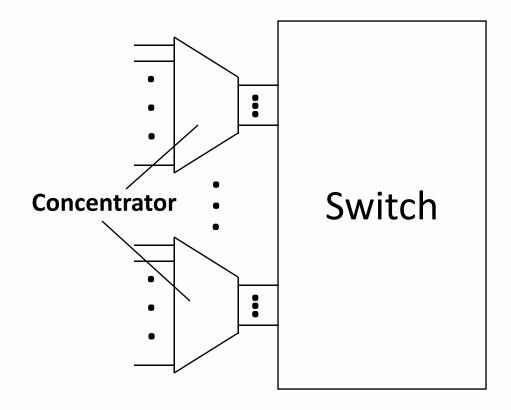
#### **Timescales of Information Transfer**

#### **Packets and Cells:**

- Transport Entities
- Variable versus fixed length
- Long versus short packets

## Messages are broken into smaller packets in network since:

- Network does not support large packet size
- Most networks are store-and-forward network, less delay under light load
- Prevent a long packet from hogging a communication channel



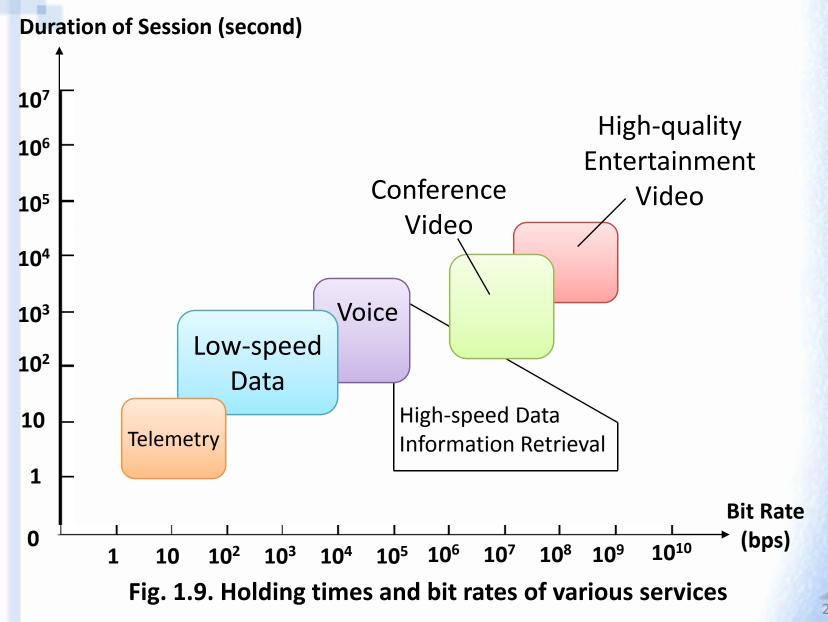
**Concentration used to reduce switching cost** 

#### **Broadband Integrated Services Network**

- Capable of supporting many different kinds of services
- Same resources can be assigned to different services at different times
- Better resources sharing (different services have different peak times)
- Support unforeseen future services

**Challenge:** How to support and control traffic of diverse characteristics in one single network?

#### **Broadband Integrated Services Network**



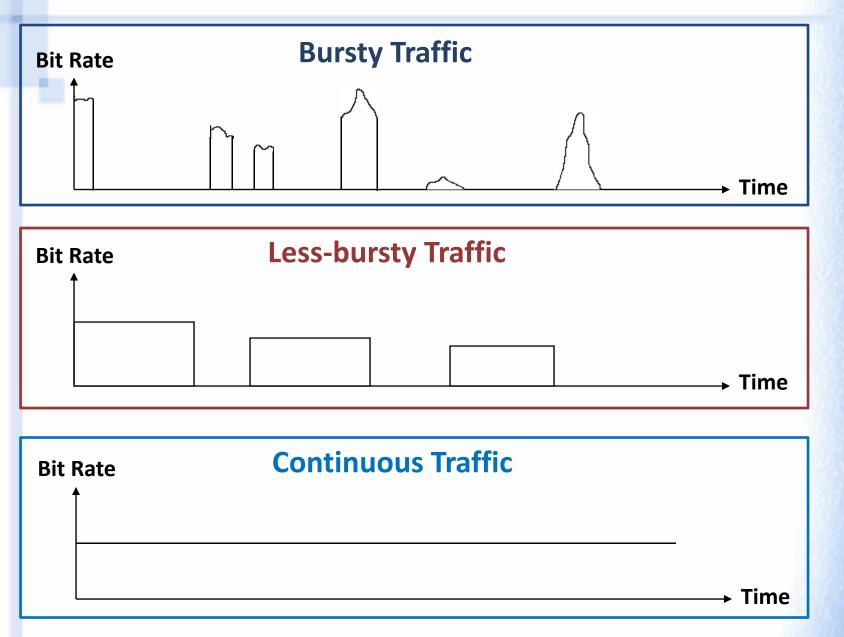


Fig. 1.10. Traffic characteristics of sources of different burstiness